

More on Functions; WHILE instructions

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- In many cases, although a block of instructions is to be repeated, it is not known before hand how many times it should be iterated.
- For example, to find an element in an array or matrix (or a word in a sequence of text), one might not have to look at **all** the elements of the array/matrix/text, since the element may be found before. In this case, the use of a FOR instruction (although possible) might not be desirable.
- In these cases, the WHILE instruction should be used , as illustrated in the next example: the **Euclid's algorithm** to find the maximum common divider between two integers.
- The WHILE instruction has the following syntax in MATLAB

while <CONDITION> WHILE-BLOCK end;

while <CONDITION> WHILE-BLOCK end;

- The behaviour of this instruction is quite intuitive. When the program reaches this instruction
	- 1. The CONDITION is assessed
	- 2. If the condition is not satisfied the WHILE-BLOCK is not executed and the program "jumps" to the next instruction.
	- 3. Otherwise, the WHILE-BLOCK is executed.
	- 4. After executing the block, the program goes back to step 1 (to assess the CONDITON again, …).
- **NOTE**: Care has to be taken in the specification of the condition and the WHILE-BLOCK. In particular, if this block does not change the variables involved in the CONDITION, so as to make it eventually false, the program **loops forever**!

Euclid's Algorithm

- The Maximum Common Divider (MCD) of two integers, can be obtained by the following algorithm.
- 1. Take the two numbers, and make them A and B, ensuring that A is no less than B.
- 2. While A is greater than B
	- Obtain C, the difference between A and B (i.e. $C = A B$);
	- Rename the numbers B and C, such that A becomes the larger of them and B the smallest.
	- Check again the condition and iterate as many times as needed.
- When one gets A equal to B, the iterations stop.
- The MCD of the initial numbers is A.

Euclid's Algorithm

Example:

• Let the numbers be 270 and 72, and see the evolution of the values of **a**, **b** and **c**.

• Hence 18 is the MCD between 270 and 72.

Euclid's Algorithm - WHILE

• The Euclid's Algorithm can be implemented with the following function:

```
function m = euclid(p, q)
% m = euclid(p, q)
% this function computes m, the maximum
% common divider between p and q.
  a = max(p,q);
  b = min(p,q);
  while a > b
     c = b – a;
     if c < b
         a = b;
% the order between these two
         b = c;
% assignments cannot change!else
        a = c; % and b remains b
     end
  end % at this point a = b
  m = b;
end
```


Euclid's Algorithm - WHILE

• A trace of the function execution shows how the values of f2, f1 and f are maintained

- We can go back to the problem referred above of finding a value in an array.
- In particular we are interested in specifying a function **find/2** that takes
	- A number as first argument; and
	- An array as second argument;

and returns

- The index of the first position where that element appears.
- **Note**: If there is no such element the function should return 0.
- Some examples:
	- **find(3, [5, 8, 4, 3, 6, 8, 2])** \rightarrow 4
	- **find(8, [5, 8, 4, 3, 6, 8, 2])** à **2**
	- **find(9, [5, 8, 4, 3, 6, 8, 2])** \rightarrow 0

- Before implementing the function we may design a convenient algorithm to solve this problem. Informally
	- While you have not found it and there is a next element
		- Look at the next element of the array to see if it is the intended one
	- Report the index of the element where you found it
- Although the skeleton of the algorithm is there, a few points must be taken care
	- 1. Where do we start from
	- 2. What if the element is not in the array
- Firstly, we must guarantee that we look at the first element, … if there is one!
- Secondly, if there are no more elements to look at, the algorithm must return 0.
- These issues may be dealt with in the specification of the **find/2** function

• The algorithm can now be implemented as function find/2, shown below

```
function k = find(v, V)% k = find(v, V)
% this function returns k, the first position, where
% v is in array V. It returns 0 if v is not present.
  found = false;
  k = 0;i = 1;n = length(V);
  while i <= n && !found % while not found and
     if v == V(i) % there is a next element to check
        k = i;found = true;
     else
        i = i + 1 ;
     end
  end
end
```


WHILE vs. FOR

- Sometimes, namely when it is known the maximum number of times a cycle might be repeated, an instruction FOR might be used to force this (max9 number of cycles
- In this case, when the condition to stop the cycle becomes True, then the cycle should be interrupted.
- In the context of a function, this may be achieved with instruction return, as below

```
function k = \text{find } 2(v, V)% k = find(v, V)
% this function returns k, the first position, where
% v is in array V. It returns 0 if v is not present.
  k = 0; % initially, the element is yet to find
   for i = 1: length(V)if v == V(i) % if the element is found in position i
        k = i; % assign the value of the function to i
        return; % and return (finish the function)
     end
  end
end
```


• A last note on the condition that could have been used in the WHILE

```
while i <= length(V) && V(i) != v
```
• As we know, trying to read an element of an array past its size reports an error

```
>> A = [ 4 7 5];
>> A(4)
    error: A(I): index out of bounds; value 4 out of bound 3
```
- Hence it is important that testing the value of the element is only done after the index is checked to be within the bound.
- In MATLAB the Boolean expression A && B (resp. A || B) is executed as follows
	- 1. Firstly, the Boolean expression A is assessed;
	- 2. If A is False (resp. True) the condition is False (resp. True)
	- 3. Otherwise B is assessed.
	- 4. The value of the condition is the value of B

Nested Functions

- As functions become more complex, their design relies on other functions, either system defined functions or user functions previously defined.
- For example if the **sin/1** function has been defined then the **tang/1** function can be defined in the obvious way.

```
function t = tan q(x)% t = \tan(x)% this function returns t, the tangent of the angle x
   s = sin(x);
  c = sqrt(1-s^2)}t = s/c;
end
```
- Hence functions can call **other** functions. Assuming the called functions terminate, the calling functions will also terminate.
- However, what happens when a function calls itself?

Recursive Functions: Factorial

- When functions call themselves, i.e. they are defined recursively, one must be careful so that they do terminate.
- Take for example the case of the function **fact/1** defined recursively to obtain the factorial of a non-negative integer (i.e the factorial/1 function, that is already predefined in MATLAB).
- This functionality could of course be defined **iteratively**, by means of the **accumulation** technique that we have seen in the previous class, implemented with a for loop.

```
function f = fact 1(n)% f = fact_1(n)
% this function returns f, the factorial of number n
  p = 1;
  for i = 1:n;
     p = p * i;
  end;
   f = p;
end
```


Recursive Functions: Factorial

• A more "mathematical" definition could however be used to guide the function implementation:

$$
n! = \begin{cases} 1 & \text{if } n < = 1 \\ n * (n-1)! & \text{if } n > 1 \end{cases}
$$

```
function f = factor(f)% f = fact (n)
% this function returns f, the factorial of number n
   if n <= 1;
      f = n;
   else
      f = n * fact(n-1);
   end
end
```
- Notice that in the implementation of this recursive function, the termination condition must be tested **before** the recursive call is made.
- Otherwise the program **loops forever**!

Recursive Functions: Factorial

• A more "mathematical" definition could however be used to guide the function implementation: $\sqrt{2}$ **1 if n <= 1**

$$
n! = \begin{cases} 1 & \text{if } n < = 1 \\ n * (n-1)! & \text{if } n > 1 \end{cases}
$$

```
function f = fact (n)
% f = fact (n)
% this function returns f, the factorial of number n
   if n <= 1;
      f = n;
   else
      f = n * fact(n-1);
   end
end
```
- Important: In the implementation of a recursive function, the termination condition is tested **before** the recursive call is made. Otherwise the program **loops forever**!
- **Note**: MATLAB has a predefined variable, **max_recursion_depth,** with a (default) value of 256, stops recursion if the depth is exceeded, thus preventing endless loops.

Recursive Functions: Maximum Common Divider

• The same recursive technique may be used to define the MCD of two numbers, taking into account that : **m if m = n** $mdc(min(m,n), abs(m-n)$ if $m \neq n$ **mdc(m,n) =**

```
function d = \text{mdc}(m, n)% d = mdc (m,n)
% this function returns d, the maximum common divider
% of integers m and n
   if m == n
      d = m;
   else
      p = min(m , n);
     q = abs(m - n);
      d = mdc(p , q);
   end
end
```
• Note again that in this recursive function, the termination condition is tested **before** the recursive call is made

• A final example of a function that is defined recursively returns the nth Fibonacci element of the series

1, 1, 2, 3, 5, 8,13, 21, 34, 55 …

- Note that in this series, every element is the sum of the two previous elements.
- Hence the function can be defined recursively as

$$
fib(n) = \begin{cases} 1 & \text{if } n <= 2 \\ fib(m-1) + fib(m-2) & \text{if } n > 2 \end{cases}
$$

- There is a (significant) difference in this case, which is the fact that the function is recursively called twice, as we will analyse later.
- But from a modelling point of view, the recursively defined function can be implemented as before.

$$
fib(n) = \begin{cases} 1 & \text{if } n <= 2 \\ fib(m-1) + fib(m-2) & \text{if } n > 2 \end{cases}
$$

```
function f = fib(n)% f = fib(n)
% this function returns f, the nth fibonnaci number
   if n \leq 2f = 1;
   else
      f = fib(n-1) + fib(n-2);
   end
end
```
- Although the termination condition is tested **before** the recursive calls are made, now there are two recursive calls and this has a big impact on the execution
- In particular, many instances of function fib, *with the same input arguments*, are called several times, in fact an **exponential** number of times!

• In fact, we can trace the computation, and see that the following calls are made

- $fib(7)$ is called 1 time
- fib(6) is called 1 times
- fib(5) is called 2 times
- fib(4) is called 3 times
- \cdot fib(3) is called 5 times
- In general,
	- fib(3) is called fib(n-2) times
	- fib(4) is called fib(n-3) times, \dots
- and fib(n) grows exponentially!

1, 1, 2, 3, 5, 8,13, 21, 34, 55, 89, 144, 233, 377, 610, 987, 1597, 2584, 4181, 6765, 10946, …

- There are two ways of avoiding this exponential explosion with double recursive functions
	- 1. use the iterative version for modelling the function
	- 2. memorize the values of the previous calls
- The iterative version, shown below, maintaining the previous 2 fibonacci numbers in two variables f2 and f1 that are added to obtain the current finonacci number.

```
function f = fibite(n)% f = fib(n)
% this function returns f, the nth fibonnaci number
% using an iterative modelling
   f = 1; f = 2 = 1; f = 1for i = 3:nf = f2 + f1;f2 = f1;f1 = f;
   end
end
```
• Note that the iterations only take place for $n \geq 3$.

• A trace of the function execution shows how the values of f2, f1 and f are maintained

- The recursive version with memorization maintains a vector as a **global** variable, i.e. a variable that is defined in the global context, and is thus visible from inside any function.
- Let us call this vector variable fib vec, and define it in the outer context (initializing the first two numbers in the fibonacci sequence to 1)

```
\geq global fib vec = zeros(1,7)
\gg fb vec(1:2) = 1
fib_vec = 1 1 0 0 0 0 0
```
- Now, any function can read from and write into this function if it identifies the variable as global, **inside** the function body.
- This is done through a global declaration, inside the function body

```
function ...
 global fib_vec;
end
```


• Now the recursive version with memorisation is easily explained.

```
If the value has not been computed yet (i.e. n > 2 & \epsilon is vec(n) \neq 0) then
   it is is computed by the (double) recursive call, and
   written in fib vec
now the value in fib_vec, can be returned
```

```
function f = fib mem(n);global fib_vec; % fib_vec identified as global
  if n > 2 && fib_vec(n) == 0 % value has yet computed
      fib vec(n) = fib mem(n-1)+fib mem(n-2);
  end
  f = fib \text{ vec}(n) = f;end;
```


Global Variables

- A last note on global variables, which have a *state* and the following life cycle.
- 1. Variables are created, in the outer context, with the declaration **global**.
- 2. Then they are assessed, either in the outer context, or within some function body.
	- a. In this case, they must be identified as global (not to be created again, only to be identified)
- 3. Eventually, they are destroyed, either because the outer context is finished, or the user wants to reset them.
	- a. In the latter case, the instruction **clear** must be used.

```
>> global vec = [ 1 2 3];
>> vec
vec = 1 2 3>> clear vec
>> vec
error: Invalid call to vec. Correct usage is:
 ...
```
Note: Some predefined variables (**pi**, **e**) are predefined global variables. If they are redefined by some assignment, they may be cleared to return to their predefined values.